COMP 5660/6660 - Evolutionary Computing - Lecture Slides

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Computational Problem Solving

• Step 1: build abstract/computational model of the real-world

¹https://quoteinvestigator.com/2011/05/13/einstein-simple/

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- Step 2: solve computationally in abstract model
- "Everything Should Be Made as Simple as Possible, But Not Simpler"¹
- Step 3: map solution back to real-world

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- A hyper-heuristic is a meta-heuristic for a space of programs

Algorithmic Toolbox

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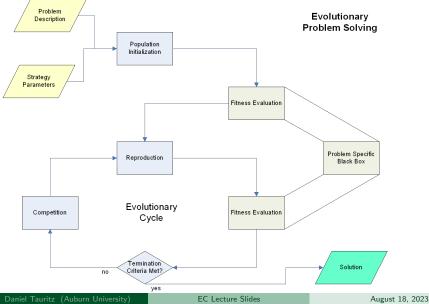
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Evolutionary Cycle



Let *F* be the decoder function from *G* (genospace) to *P* (phenospace) and x^* be the global optimum.

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- If F is not surjective and x^{*} ∉ F(G), then the EA cannot find the global optimum. Therefore one should think twice before choosing a non-surjective decoder function if one cannot guarantee that the global optimum is still reachable.
- F does not need to be injective, but realize there is less to search if F is injective so there should be sufficient compensation, such as limiting F(G) to valid solutions in a constraint satisfaction problem.

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- fitness $(p) = \sum_{i=1}^{n} (v_i \cdot g_i)$
- **2** Modify fitness(p) to exclude items that would exceed C_{max}